

Proving versus testing in climate impact research

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PIK: Potsdam Institute for Climate Impact Research

“At PIK researchers in the natural and social sciences work together to study global change and its impacts on ecological, economic and social systems. They examine the Earth system’s capacity for withstanding human interventions and devise strategies for a sustainable development of humankind and nature.

PIK research projects are interdisciplinary and undertaken by scientists from the following Research Domains: Earth System Analysis, Climate Impacts and Vulnerabilities, Sustainable Solutions and Transdisciplinary Concepts and Methods.”

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Brave new world.

Lots of scope for conceptual analysis.

Examples:

vulnerability

resilience

adaptive capacity

mitigation

sustainability

...

Best tool for the job

Opinions differ...

physicist: partial differential equations

economist: utility functions

social scientist: subject interviews

mathematician: sets, lattices, categories

Example: “Vulnerability”

“... a human condition or process resulting from physical, social and environmental factors which determine the likelihood and damage from the impact of a given hazard” (UNDP Annual Report, 2004)

“Vulnerability [...] is a way of conceptualizing what may happen to an identifiable population under conditions of particular risk and hazards.” (Cannon et al. 2004)

“... the degree to which a system is susceptible to and unable to cope with, adverse effects of climate change, including climate variability and extremes. ” (The Intergovernmental Panel on Climate Change, 2007)

Vulnerability formalization

Basic elements:

S -- set of states
 $Trj = List\ S\ or\ Trj = T \rightarrow S$ -- trajectory
 $h : Trj \rightarrow H$ -- harm along a trajectory

Vulnerability formalization

“Possible” future:

$[trj_1, \dots, trj_n]$ scenarios

or

$[(trj_1, p_1), \dots, (trj_n, p_n)]$ stochastic uncertainty

or

fuzzy, finitely additive probabilities, . . .

or

combinations thereof

Vulnerability formalization

“Possible” future: $F \text{ Trj}$ for a functor F

$$p : \text{State} \rightarrow F \text{ Trj}$$

$F h \circ p : \text{State} \rightarrow F H$ possible future harm

Vulnerability is a measure of this possible future harm:

$$m : F H \rightarrow V \text{ for some preorder } V$$

Vulnerability formalization

$$v : S \rightarrow V$$

$$v = m \circ F h \circ p$$

In Haskell, using “user-friendly” names:

$$vulnerability = measure \circ fmap harm \circ possible$$

Formalization as DSL creation.

Monotonicity condition

Monotonicity condition for vulnerability measures:

For all $f : H \rightarrow H$ such that f is increasing, for all $fh : F H$, we have

$$m fh \sqsubseteq m (F f fh)$$

or, if \sqsubseteq is partial

$$\neg (m (F f fh) \sqsubseteq m fh) \quad \text{-- for suitably defined } \sqsubseteq$$

If every possible harm is increased, the total measure of possible harm should not be decreased.

Monotonicity condition examples

$H = \text{Nat}, F = \text{List}$

1. $V = \text{Nat}, m = \text{maximum}$ works
2. $V = \text{Real}, m = \text{average}$ works
3. $V = \text{Nat}, m = \text{most frequent value}$ fails

Monotonicity condition examples

$$H = \text{Nat}, F = \text{SP}$$

1. $V = \text{Nat}, m = \text{maximum value}$ works
2. $V = \text{Real}, m = \text{expected value}$ works
3. $V = \text{Nat}, m = \text{likeliest value}$ fails

Testing the monotonicity condition

Testing the monotonicity condition using QuickCheck:

$$\text{testMonotonicity measure geninc fh} = \text{forAll geninc} \\ (\lambda \text{inc} \rightarrow ((\text{measure fh}) \sqsubseteq (\text{measure (fmap inc fh)})))$$

Problems:

we need to create a custom generator for every harm type (and every preorder on the harm type)

pretty bad coverage for high dimensions

Testing the monotonicity condition

Testing the monotonicity condition:

can fail because of overflow or round-off errors

hard to distinguish between “conceptual” and “implementational” errors

Proving the monotonicity condition

Formulating the monotonicity condition in Agda is like writing the tests in QuickCheck.

record Preorder (A : Set) : Set **where**

field

le : A → A → Bool

LE : A → A → Set

LE a₁ a₂ = lift (le a₁ a₂)

field

reflexive : (a : A) → LE a a

transitive : (a₁ a₂ a₃ : A) →

LE a₁ a₂ → LE a₂ a₃ → LE a₁ a₃

Proving the monotonicity condition

record Functor ($F : Set \rightarrow Set$) : *Set1* **where**

field

$fmap : \{ A B : Set \} \rightarrow (A \rightarrow B) \rightarrow (F A \rightarrow F B)$

....

$LE : \{ A : Set \} \rightarrow (P : Preorder A) \rightarrow (a_1 a_2 : A) \rightarrow Set$

$LE P a_1 a_2 = lift (le a_1 a_2)$ **where** $le = Preorder.le P$

record Increasing ($A : Set$) ($P : Preorder A$) ($f : A \rightarrow A$) : *Set* **where**

field

$increasing : (a : A) \rightarrow LE P a (f a)$

Proving the monotonicity condition

```
record VulnMeasure (H V : Set) (PH : Preorder H) (PV : Preorder V)  
  (F : Set → Set) (Func : Functor F) (m : F H → V) : Set where  
  fmap = Functor.fmap Func  
  field  
    mon : (inc : H → H) → (Increasing H PH inc) → (fh : F H) →  
      LE PV (m fh) (m (fmap inc fh))
```

Proving the monotonicity condition

It is easy to construct a term of type

VulnMeasure H H P P List ListFunc maximum

for any harm type H .

Similarly for *average*, *expected*, etc. (with appropriate assumptions on H).

In constructing these proofs, some properties will be *postulated* (e.g. associativity of addition on H) which do not hold, e.g. of *Float*.

Compatibility condition

Usually, different models for possible trajectories of the same system have the same state space S and the same trajectory space Trj , but differ in the choice of F , e.g. some give possible scenarios ($List\ Trj$), others stochastic information ($SP\ Trj$).

When are two vulnerability assessments with models of different types “compatible” ?

The harm evaluation stays the same ($h : Trj \rightarrow H$), but the domain of the vulnerability measure $m : F\ H \rightarrow V$ has to change.

A related problem

Intuitively speaking, $m_1 : F H \rightarrow V$ and $m_2 : F H \rightarrow V$ are compatible if they rank things in the same way:

$$\forall fh_1 fh_2 : F H \\ m_1 fh_1 \sqsubseteq m_1 fh_2 \equiv m_2 fh_1 \sqsubseteq m_2 fh_2$$

i.e. if the induced preorders on $F H$ are order-isomorphic.

Compatibility condition

We can “reuse” this idea if one of the representations of “possible” can be embedded in the other, i.e. if there exists an *injective* natural transformation from one to the other.

Let F_1, F_2 be two functors, $\tau : F_1 \rightarrow F_2$ an injective natural transformation, H a set, \sqsubseteq_1 and \sqsubseteq_2 preorders on $F_1 H$ and $F_2 H$ respectively. Then \sqsubseteq_1 and \sqsubseteq_2 are compatible with respect to τ iff

$$\forall x_1 x_2 : F_1 H \\ x_1 \sqsubseteq_1 x_2 \equiv \tau x_1 \sqsubseteq_2 \tau x_2$$

Conclusions

The ability to easily formulate and prove high-level conditions is another advantage of using a (dependently typed) functional programming language for formalization of concepts.

Ideally, it would be as easy to prove these kind of properties as it is to implement QuickCheck tests.

But...

... we're not there yet:

1. We lack a good enough tutorial
2. Agda standard library is intimidating
3. Not enough experience with proof reuse